## Translation and rotation in the plane



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Material taken form the book: J. C. Latombe, Robot motion planning.

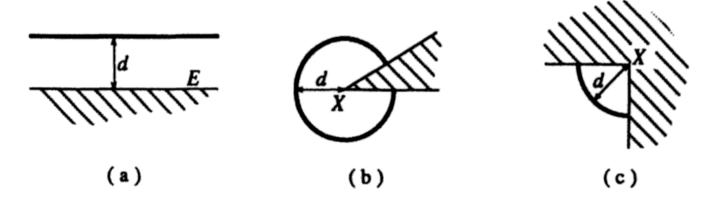
## The setup

- The working space is  $W = \mathbb{R}^2$ .
- The robot *A* is line segment of length *d*, with endpoints *P* and *Q*. The point *P* is located at (x, y), the orientation is taken from an angle  $\theta \in [0, 2\pi)$ . Thus the configuration space *C* of *A* is  $\mathbb{R}^2 \times S^1$ , each configuration is  $(x, y, \theta)$ .
- The union of all obstacles form a polygonal region B.
- *CB* is a three dimensional region in  $\mathbb{R}^2 \times [0, 2\pi)$ .

## **Critical curves**

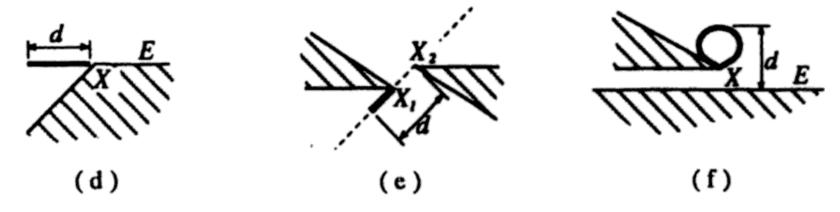
These curves can be expresed with an algebraic formula.

Obstacle edges are said to be critical curves of type 0.



a) The line segment at distance d from an obstacle edge E. This line has the same length as E. E is a critical curve of **type 1**.

b) and c) The arc of radius d, centered at the obstacle vertex X and bounded the two lines that form the vertex. X is a critical curve of **type 2**.



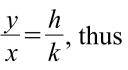
d) Let *E* be an obstacle edge and *X* a convex vertex at one endpoint of *E*. *PQ* is contained in the line containing *E*. The line traced as *A* slides and *Q* is in *E* is a critical curve of **type 3**.

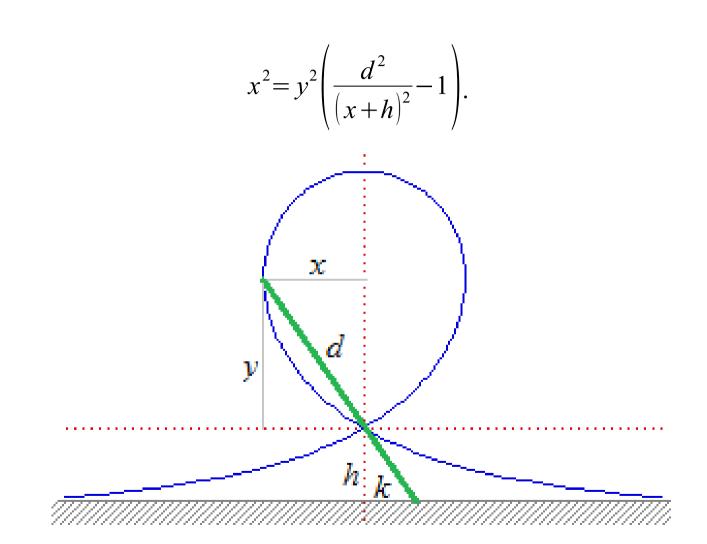
e)  $X_1$  and  $X_2$  are two convecx obstacle vectices, A is bitangent to them. The line traced by P as A moves is a critical curve of **type 4**.

f) *E* is and obstacle edge, *X* is a convex obstacle vertex that is not and endpoint of *E*. The distance between *E* and *X* is *h*. If h < d, the curve traced by *P*, as *A* moves in a way such that it touches both *E* and *X* is a critical curve of **type 5**.

This curve is a conchoid of Nicodemes, with

 $d^{2} = (y+h)^{2} + (x+k)^{2},$ 





The set of critical curves is finite.

Every critical curve is a smooth curve algebraic curve of degree 1 (0, 1, 3, 4), 2 (2), or 4 (5).

The configuration space will be divided using this critical curves using a connectivity graph G.

G is an undirected graph, each node is a cell in C. Two nodes are connected if the cells that they represent are adjacents.

A position (x, y) is admisible if there exists at least one orientation  $\theta$  such that  $(x, y, \theta) \in C_{\text{free}}$ .

A non-critical region *R* is defined in base of a contact of *A* with obstacles, it is defined as  $F(x, y) = \{\theta | (x, y, \theta) \in C_{\text{free}} \}.$ 

If *A* is free for all  $\theta$ , then

$$F(x, y) = [0, 2\pi),$$

else

$$F(x, y) = a$$
 finite set of intervals

For each maximally connected interval  $(\theta_1, \theta_2) \subseteq F(x, y)$ , let  $s_1$  and  $s_2$  the contacts by  $A(x, y, \theta_1)$ and  $A(x, y, \theta_2)$ . The contacts could be on either a vertex X or an edge E.

Each interval has associated two contacts, one is a clockwise (at  $\theta_c$ ) the other is counterclockwise (at  $\theta'_c$ ).

Let  $\sigma(x, y)$  be the set of all the pairs  $[s(x, y, \theta_c), s(x, y, \theta'_c)]$ , the interval  $(\theta_c, \theta'_c) \subseteq F(x, y)$ .

If  $F(x, y) = [0, 2\pi)$ , we write  $\sigma(x, y) = \{[\Omega, \Omega]\}$ .

Given a pair  $[s_1, s_2] \neq [\Omega, \Omega]$ , we denote an unique orientation  $\lambda(x, y, s_1)$  such that  $A(x, y, \lambda(x, y, s_1))$  touches the clockwise stop  $s_1$ .

Let *R* be a non-critical region.

A cell is defined as

$$cell(R, s_1, s_2) = \left\{ (x, y, \theta) \middle| (x, y) \in R \land \theta \in (\lambda(x, y, s_1), \lambda(x, y, s_2)) \right\}$$

Two cells  $\kappa = cell(R, s_1, s_2)$  and  $\kappa' = cell(R', s'_1, s'_2)$  are adjacent if and only if:

• The boundaries  $\partial R$  and  $\partial R'$  of R and R' share a critical curve section  $\beta$ , and

• 
$$\forall (x, y) \in int(\beta) : (\lambda(x, y, s_1), \lambda(x, y, s_2)) \cap (\lambda(x, y, s'_1), \lambda(x, y, s'_2)) \neq \emptyset$$