

# Recovering trees from the cohomology ring of their configuration spaces

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## Abstract

Given a tree  $T$ , the cohomology ring of its unordered configuration space  $H^*(UD^nT)$  is an exterior face algebra if  $T$  is a binary core tree (if by removing the leaves from  $T$  we obtain a binary tree), or if  $n = 4$ . This means that every cup product is determined by a simplicial complex  $K_nT$ . In this paper we show how to recover the tree  $T$  from the simplicial complex  $K_nT$  when  $n = 4$ .

## 1 Introduction

For a finite graph  $G$  and a positive integer  $n$ , the discretized unlabelled configuration space on  $n$  points of  $G$  is defined as

$$UD^nG = \{\{x_1, \dots, x_n\} : x_i \in V(G) \cup E(G), x_i \cap x_j = \emptyset \text{ if } i \neq j\}.$$

Configuration spaces of graphs have been widely studied, in particular, their cohomology ring is well known when  $G$  is a tree. Using discrete Morse theory techniques, D. Farley gave in [2] an efficient description of the additive structure of the cohomology ring of  $UD^nG$  when  $G$  is a tree  $T$ . Later, and in order to get to the multiplicative structure, the Morse theoretic methods were replaced in [3] by the use of a Salvetti complex  $\mathcal{S}$  obtained by identifying opposite faces of cells in  $UD^nT$ . Being a union of tori,  $\mathcal{S}$  has a well understood cohomology ring.

Given a tree  $T$  we can construct another tree  $F(T)$ , where the vertex set of  $F(T)$  is  $V(F(T)) = \{x \in V(T) : d(x) > 2\}$  and two vertices are adjacent in  $F(T)$  if the unique path joining them in  $T$  does not contain any other vertex of  $F(T)$ . We say that a tree  $T$  is a binary core tree if  $F(T)$  is a binary tree. In [1], the following theorem was proven:

**Theorem 1.1.** [1] *Given an integer  $n \geq 4$ , the cohomology ring  $H^*(UD^nT)$  is an exterior face ring if either  $n = 4$ , or  $T$  is a binary core tree.*

This means that when  $T$  is a binary core tree or  $n = 4$ , all products in  $H^*(UD^nT)$  are given by a simplicial complex  $K_nT$ , where the vertices of  $K_nT$  are the basis elements of

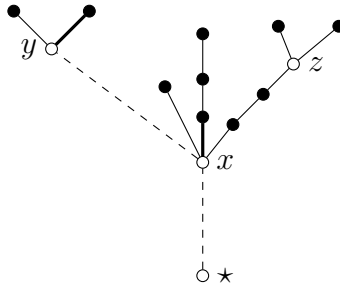
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dimension 1, and a set of  $k$  vertices forms a simplex in  $K_n T$  if and only if the cup product of the corresponding  $k$  elements is non zero. The simplicial complex  $K_n T$  is called the  $n$ -interaction complex of  $T$ .

The 1-skeleton of the space  $UD^n T$ , coincides precisely with the token graph  $F_n(T)$  of  $T$  on  $n$  tokens. In [4], Fabila-Monroy and Trujillo-Negrete proved that if a graph does not have induced cycles of length four, or induced diamonds (a graph isomorphic to  $K_4 \setminus \{e\}$  the complete graph on four vertices without an edge), then the graph  $G$  is uniquely reconstructible from  $F_k(G)$  and gave an algorithm to do so in polynomial time. One then might ask if there is a topological version of this, this is, if we can recover the tree  $T$  from a topological property of  $UD^n T$ . We are going to prove that we can recover the tree  $T$  from the  $n$ -interaction complex  $K_n T$  in the case when  $n = 4$ , for any tree  $T$ .

We shall assume that the tree  $T$  is embedded in the plane and has as root a vertex of degree one  $\star$ . The edges inciding in a vertex  $x$  are enumerated by this embedding and we fix the edge that lies on the unique  $x\star$  path to be the 0th edge. We shall say that a vertex  $y$  lies on  $x$ -direction  $i$  for  $1 \leq i \leq d(x) - 1$  if  $x$  belongs to the unique path joining  $y$  and the root vertex, and if this path contains the  $i$ th edge inciding in  $x$ . This also implies that  $x$  lies on  $y$ -direction 0. If  $x$  does not belong to the  $y\star$  path and  $y$  does not belong to the  $x\star$  path we shall also say that  $x$  lies on  $y$ -direction 0 and  $y$  lies on  $x$ -direction 0, in which case we say that the vertices  $x$  and  $y$  are not stacked. Given two vertices of degree at least three  $x$  and  $y$ , we shall say that  $x < y$  if they are stacked and  $x$  lies on  $y$ -direction zero, or if they are not stacked and there exists  $z$  such that  $x$  lies on  $z$ -direction  $i$  and  $y$  lies on  $z$ -direction  $j$  with  $i < j$ .



**Figure 1:** The tree  $T$  containing the vertices  $x, y$  and  $z$ .

**Example 1.2.** Consider the tree shown in Figure 1, where the dashed edges represent paths of any length greater than two. Then the vertices  $x$  and  $y$  are stacked, the vertices  $z$  and  $x$  are also stacked but the vertices  $y$  and  $z$  are not stacked.

We are going to define the interaction complex  $K_n T$  only for the case when  $n = 4$ , in which case it is a graph. For the the definition of the interaction complex in the general case we refer the reader to [1].

## 2 The graph $K_4T$

**Definition 2.1.** *The graph  $K_4T$  is defined as follows:*

- *The vertices are 4-tuples  $(k, x, p, q)$  such that  $x$  is a vertex of  $T$  with  $d(x) \geq 3$ ,  $k$  is a non-negative integer and  $p$  and  $q$  are integer vectors having non-negative entries such that  $p$  has at least one positive entry, the sum of their lengths  $l(p) + l(q)$  is  $d(x) - 1$  and the sum of their entries is  $3 - k$ .*
- *Let  $v = (k_1, x_1, p_1, q_1)$  and  $w = (k_2, x_2, p_2, q_2)$  be two vertices, then*
  1. *if the vertices  $x_1$  and  $x_2$  are not stacked, then  $(v, w) \in E(K_4T)$  if  $k_1 + k_2 \geq 4$*
  2. *if  $x_2$  lies on  $x_1$ -direction  $i$  with  $l(p_1) \geq i$ , then  $(v, w) \in E(K_4T)$  if  $p_{1,i} > 4 - k_2$  or  $p_{1,i} + k_2 = 4$  and there exists a  $j \neq i$  such that  $p_{1,j} \neq 0$*
  3. *if  $x_2$  lies on  $x_1$ -direction  $i$  with  $l(p_1) < i$ , then  $(v, w) \in E(K_4T)$  if  $q_{1,i-l(p_1)} > 4 - k_2$*

In particular, the graph  $K_4T$  has a lot of isolated vertices, as we can see in the following lemma.

**Lemma 2.2.** *A vertex  $v = (k, x, p, q)$  is an isolated vertex if one of the following conditions hold:*

- $k = 1$
- *the sum of the entries of  $p$  is two.*

*Proof.* Assume first that  $k = 1$ . Notice that since  $l \in \{0, 1, 2\}$ , if a vertex  $w = (l, y, r, s)$  is such that  $x$  and  $y$  are not stacked, then both  $k$  and  $l$  must be two for  $v$  and  $w$  to be adjacent. If  $x$  and  $y$  are stacked, every entry of  $p = (p_1, \dots, p_{l(p)})$  is at most two, thus  $2 \geq p_i$  and  $4 - l \geq 2$ . Finally notice that every entry of  $q = (q_1, \dots, q_{l(q)})$  is at most one so  $1 \geq q_{i-l(p_1)}$  and  $4 - k_2 \geq 2$ . This proves that every vertex with  $k = 1$  is isolated. Assume now that the sum of the entries of  $p$  is two, and consider again  $w = (l, y, r, s)$ . This means that  $k = 0$  and  $q$  has one unique entry with value one. Then again  $2 \geq p_i$ ,  $4 - l \geq 2$ ,  $1 \geq q_{i-l(p_1)}$  and  $4 - k_2 \geq 2$  thus  $v$  is an isolated vertex.  $\square$

For our purposes, the isolated vertices are not relevant, so we shall ignore them.

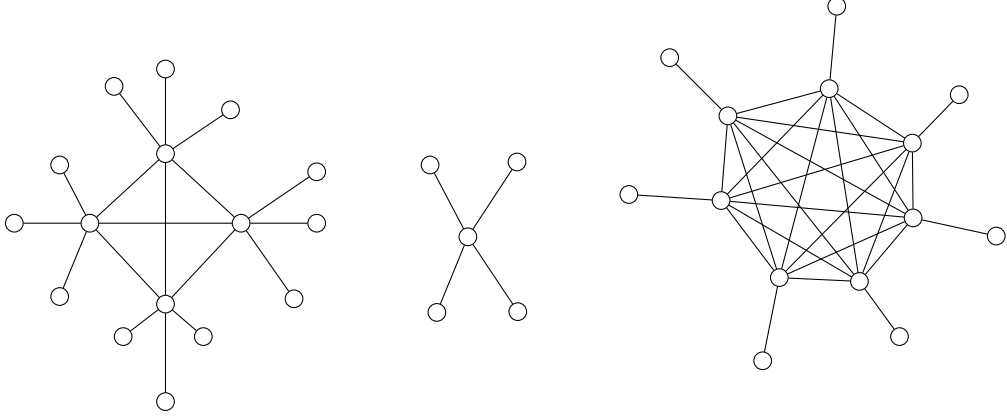
**Proposition 2.3.** *Given a vertex  $x$  of  $T$ , fix  $i \in \{1, \dots, d(x) - 1\}$  and let*

$$\Lambda_i = \{(k, x, p, q) \in K_4T : p_i \geq 2 \text{ if } i \leq l(p) \text{ or } q_{i-l(p)} \geq 2 \text{ if } i > l(p)\}.$$

*Then the set  $\Lambda_i$  has cardinality  $\frac{1}{2}(d(x) - 1)(d(x) - 2)$ , and this is also the amount of vertices having  $k = 2$ .*

*Proof.* Assume first that  $k = 0$ . Since the sum of the entries of  $p = (p_1, \dots, p_{l(p)})$  and  $q$  must be at most  $3 - k = 3$ , and one entry is fixed to be 2, there can only be one additional non zero entry. Recall also that  $p$  must have at least one zero entry, and since we are considering only vertices in which the sum of the entries of  $p$  is not two,  $p$  must have an

additional non zero entry or  $p_i = 3$ . Then for every possible length of  $p$ , we have  $l(p)$  options, one for each additional 1 as an entry. Since  $l(p)$  varies between 1 and  $d(x) - 2$  we have  $\sum_{i=1}^{d(x)-2} i = (d(x) - 1)(d(x) - 2)$ . The same happens if  $k = 2$ , since the remaining 1 must be an entry of  $p$ .  $\square$



**Figure 2:** The graph  $K_4T$ , the tree  $T = K_{1,4}$  and an example of the graph  $KP_7$ .

**Example 2.4.** When  $T = K_{1,m}$ , the graph  $K_4T$  is the graph obtained from the complete graph  $K_m$  by adding to each vertex,  $\frac{1}{2}(m - 1)(m - 2)$  new neighbours as in Figure 2.

The following proposition follows from the definition of  $K_4T$ .

**Proposition 2.5.** Let  $v = (k, x, p, q)$  be a vertex in  $K_4T$ . Then the degree of  $v$  is the amount of vertices in  $F(T)$  which lie on direction  $i$ , where

- $i = 0$  if  $k = 2$ ,
- $i$  is such that the  $i$ th entry of  $p$  is at least two if  $i \leq l(p)$ ,
- $i$  is such that the  $(i + l(p))$ th entry of  $q$  is two if  $i > l(p)$ .

**Corollary 2.6.** Given a binary tree  $T$ , the number of leaves in  $K_4T$  is the number of leaves in  $F(T)$ .

*Proof.* Let  $w$  be a leaf in  $F(T)$  different from the root, and let  $v$  be its unique neighbour. Assume  $w$  lies on  $v$ -direction  $i$ . Then the vertex  $u$  is a leaf in  $K_4T$  and has as unique neighbour the vertex  $(2, v, 1, 0)$ , where  $u = (0, w, p, q)$  is such that  $u = \begin{cases} (0, w, 3, 0) & \text{if } i = 1 \\ (0, w, 1, 2) & \text{if } i = 2. \end{cases}$  If  $w$  is the root vertex with unique neighbour  $v$ , then in  $K_4T$  the vertex  $(2, v, 1, 0)$  is a leaf and its unique neighbour is  $(0, w, 3, 0)$ .  $\square$

To simplify notation, we shall often write 0 to denote the vector having every entry zero. Let  $L(G)$  denote the set of leaves of a graph  $G$ .

**Proposition 2.7.** If two vertices  $u$  and  $v$  in  $K_4T$  have the same neighbourhood then  $u = (k, x, p, q)$  and  $v = (l, x, r, s)$ .

*Proof.* Assume  $u = (k, x, p, q)$  and  $v = (l, y, r, s)$  have the same (non empty) neighbourhood for  $x \neq y$ . Notice first that  $u$  and  $v$  are not adjacent since  $K_4T$  has no loops. Notice also that  $u$  and  $v$  can not be both leaves, since that would mean that  $u = (2, x, p, 0)$  and  $v = (2, y, r, 0)$  thus  $(u, v) \in E(K_4T)$ , a contradiction.

Assume first that  $x$  and  $y$  are not stacked. This means that either  $k$  or  $l$  is different than 2. Assume without loss of generality that  $k \neq 2$ , thus  $k = 0$ . Since  $N(u) \neq \emptyset$ , there exists  $w \in N(u)$ ,  $w = (2, z, p', q')$  and  $z$  lies on  $x$ -direction  $i$ . This is a contradiction since  $w \in N(v)$  and the vertices  $x$  and  $y$  are not stacked.

Now we may assume that the vertices  $x$  and  $y$  are stacked, and assume without loss of generality that  $y$  lies on  $x$ -direction  $i$ . This means that  $k = 0$  and  $l = 2$ . If  $x$  is not a leaf in  $F(T)$ , the root vertex  $z$  of  $F(T)$  is such that  $(0, z, p', q')$  is adjacent to  $v$  for some vectors  $p'$  and  $q'$ . But this means that  $(0, z, p', q')$  is also adjacent to  $u = (0, x, p, q)$  which is a contradiction. If  $x$  is a leaf, since we are assuming that  $y$  lies on  $x$ -direction  $i$ ,  $x$  must be the root vertex and  $u = (0, x, p, 0)$ . Then  $v = (2, y, r, 0)$ , and  $w = (2, z, p', 0) \in N(u)$  with  $z$  the unique neighbour of  $x$ . Since  $z$  and  $y$  are stacked,  $w \notin N(v)$  which is a contradiction. Hence  $x = y$ .  $\square$

**Corollary 2.8.** *If  $T$  is a binary tree, then there are no two vertices having the same non empty neighbourhood in  $K_4T$ .*

*Proof.* Assume  $v$  and  $w$  have the same neighbourhood. By Proposition 2.7, we must have that  $v, w \in \{(0, x, 1, 2), (0, x, 3, 0), (2, x, 1, 0)\}$ . If  $v = (2, x, 1, 0)$  and  $w \neq v$ , then for every vertex  $y$  such that  $y$  lies on  $x$ -direction 0, we have that if  $x$  lies on  $y$ -direction  $i$ , then  $u \in N(v)$  where  $u = \begin{cases} (2, y, 1, 0) & \text{if } i = 0 \\ (0, y, 3, 0) & \text{if } i = 1 \\ (0, y, 1, 2) & \text{if } i = 2. \end{cases}$  But, since  $w \neq v$ , we have that  $u$  can not be

adjacent to  $w$ . This means that there are no vertices lying on  $x$ -direction 0, but this implies that  $N(v) = \emptyset$ . So assume that  $v = (0, x, 1, 2)$  and  $w = (0, x, 3, 0)$ . Then again, if there exist vertices  $y, z \in T$  such that  $y$  lies on  $x$ -direction 1 and  $z$  lies on  $x$ -direction two, the vertex  $(2, y, 1, 0)$  is adjacent to  $v$  but not to  $w$  and the vertex  $(2, z, 1, 0)$  is adjacent to  $w$  but not to  $v$ . This means that  $x$  is a leaf in  $F(T)$ , but this implies that  $N(v) = \emptyset = N(w)$ .  $\square$

### 3 Recovering the tree $T$

We shall first consider the case when  $T$  is a binary tree before considering the general case, since this case is easier both notation-wise and mathematically.

#### 3.1 Binary trees

Throughout this section we shall assume that  $T$  is a binary tree. The main reason for this is Proposition 2.7. If  $T$  is a binary tree, then  $T$  is completely determined by  $F(T)$ .

**Proposition 3.1.** *Assume  $T$  is a binary tree and let  $K_m$  be a complete subgraph of  $K_4T$  for  $m \geq 3$ . Then in  $F(T)$  there exists an independent set  $I$  of  $m$  vertices.*

*Proof.* Let  $v_i = (k_i, x_i, p_i, q_i)$  for  $1 \leq i \leq m$  be the vertices of  $K_m$  in  $K_4T$ . Assume that  $x_j$  lies on  $x_i$ -direction  $t \in \{1, 2\}$  for some  $i \neq j$ . Then  $v_j = (2, x_j, 1, 0)$  and  $v_i = (0, x_i, 3, 0)$

or  $v_i = (0, x_i, 1, 2)$  (if  $t = 1$  or  $t = 2$  respectively). Since  $v_i$  multiplies every vertex  $v_l$ , this means that  $x_l$  lies on  $x_i$ -direction  $t$  for every  $l \in \{1, \dots, m\} \setminus \{i\}$ . Recall that  $v_j$ , and hence every  $v_l$ , must be of the form  $(2, x_l, 1, 0)$  for  $l \in \{1, \dots, m\} \setminus \{i\}$ . This means that  $x_l$  lies on  $x_j$ -direction 0 and vice versa  $x_j$  lies on  $x_l$ -direction 0 for  $j, l \in \{1, \dots, m\} \setminus \{i\}$ . This immediately implies that  $\{x_1, \dots, x_m\} \setminus \{x_i\}$  is an independent set of  $m - 1$  vertices. Now assume the edge  $(x_i, x_j)$  exists in  $F(T)$  for some  $j \in \{1, \dots, m\} \setminus \{i\}$ . Then  $x_j$  lies on  $x_i$ -direction 0 but this implies that  $x_i$  also lies on  $x_l$ -direction for  $l \in \{1, \dots, m\} \setminus \{i, j\}$  which is a contradiction. Hence  $I = \{x_1, \dots, x_m\}$  is an independent set in  $F(T)$ .  $\square$

**Corollary 3.2.** *Assume that  $T$  is a binary tree, and that  $F(T)$  has  $m$  leaves. Then there exists in  $K_4T$  a subgraph isomorphic to  $K_m$  such that each vertex in  $K_m$  is adjacent to exactly one leaf.*

Let  $KP_m$  be the subgraph obtained by attaching a leaf to each vertex of a complete subgraph  $K_m$  (see Figure 2 (right)).

Let  $Kp$  denote the subgraph of  $K_4T$  consisting of the complete subgraph mentioned in Corollary 3.2 together with the adjacent leaves. If we delete from  $K_4T$  the subgraph  $Kp$ , we obtain the subgraph of  $K_4T$  consisting only of vertices  $(k, x, p, q)$  such that  $x$  is not a leaf of  $F(T)$ . Moreover, the amount of leaves in  $K_4T - Kp$  corresponds to the amount of leaves in  $F(T) - L(F(T))$  and we can apply Corollary 3.2 again.

This means that if  $S_1$  be the subgraph of  $K_4T$  isomorphic to  $KP_{m_1}$  for some  $m_1$ , then  $K_4T - S_1$  contains  $S_2$ , a subgraph isomorphic to  $KP_{m_2}$ . We can continue this way until  $K_4T - \cup_{i=1}^n S_i$  is isomorphic to either a  $KP_3$ , a  $KP_2 = P_4$  or a  $KP_1 = K_2$  (since  $T$  is binary), and notice that  $m_1 \geq m_2 \geq \dots \geq m_n$ .

**Definition 3.3.** *Let  $V(K_4T) = \cup_{i=1}^n S_i$  such that the induced subgraph generated by  $S_i$  is isomorphic to  $KP_{m_i}$  as in the previous discussion. We say that a vertex  $v \in V(K_4T)$  belongs to level  $i$  if it is a vertex of  $S_i$ . Moreover,  $V(S_i) = L_i \cup R_i$  where  $L_i$  is the set of leaves of  $S_i$ .*

We are now ready to recover the tree  $T$  from  $K_4T$  when  $T$  is a binary tree. First, notice that  $V(K_4T) = \cup_{i=1}^n V(S_i)$  and since each  $S_i$  has  $2m_i$  vertices, we are going to label each vertex of  $L_i$  and  $R_i$  with the labels  $\{x_1^i, \dots, x_{m_i}^i\}$  and  $\{y_1^i, \dots, y_{m_i}^i\}$  respectively, thus in  $S_i$ ,  $y_j^i$  is the only neighbour of  $x_j^i$ . We are going to construct a tree  $F_T$  isomorphic to  $F(T)$ .

**Definition 3.4.** *The tree  $F_T$  is defined as follows:*

- *The vertices of  $F_T$  are  $V(F_T) = z_0 \cup \left( \bigcup_{i=1}^n \bigcup_{j=1}^{m_i} z_j^i \right)$ , and*
- *$N(z_0) = \bigcup_{j=1}^{m_n} z_j^n$*
- *$(z_j^i, z_l^k) \in E(F_T)$  if  $k = \min\{m > i : (x_l^m, y_j^i) \in E(K_4T)\}$ .*

We now must prove that this tree is indeed isomorphic to  $F(T)$ .

**Lemma 3.5.** *The trees  $F(T)$  and  $F_T$  are isomorphic.*

*Proof.* We proceed by induction over  $n$ , and assume  $n = 1$ . This means that  $K_4T = S_1 \cong KP_{m_1}$  which means that  $F(T)$  is a star  $F(T) \cong K_{1,m_1}$ . On the other hand,  $F_T$  consists only of the vertices  $z_0$  together with  $\cup_{i=1}^{m_1} z_i^1$ , but since  $n = 1$  we have that every vertex is adjacent to  $z_0$ , thus  $F_T = K_{1,m_1}$ .

Assume now, that the lemma is valid for values smaller than  $n$ . We have that  $K_4T \setminus S_1 = K_4T'$  where  $T'$  is such that  $F(T')$  is obtained from  $F(T)$  by removing every leaf. Then  $V(K_4T') = \cup_{i=2}^n S_i$  thus by induction hypothesis,  $F_{T'} = F(T') = F(T) \setminus L(F(T))$ . Notice that the amount of vertices in  $F(T) \setminus F(T')$  is the amount of leaves in  $F(T)$  which is  $m_1$ , which is also the amount of vertices in  $F_T \setminus F_{T'}$ . Take a vertex  $w \in F(T)$ , and let  $v \in F(T)$  be a leaf adjacent to  $w$ . By Corollary 2.6, one of the following holds:

- the vertex  $(0, v, 1, 2)$  is a leaf in  $K_4T$  and  $(2, w, 1, 0)$  is its unique neighbour,
- the vertex  $(0, v, 3, 0)$  is a leaf in  $K_4T$  and  $(2, w, 1, 0)$  is its unique neighbour, or
- the vertex  $w$  is the root vertex of  $T$ , the vertex  $(2, v, 1, 0)$  is a leaf in  $K_4T$ , and its unique neighbour is  $(0, w, 3, 0)$ .

Assume the first case holds. Then  $(0, v, 1, 2) = x_j^1$  and  $(2, w, 1, 0) = y_i^1$  for  $1 \leq i, j \leq m_1$ . Since  $v$  is not a leaf in  $F(T)$ , there exists  $u \in V(F(T))$  such that  $(v, u) \in E(F(T))$  and  $u$  lies on  $v$ -direction 0. This means that  $(2, v, 1, 0)$  is adjacent to  $(0, u, 3, 0)$  or  $(0, u, 1, 2)$ . We can assume without loss of generality that  $(0, u, 3, 0)$  is adjacent to  $(2, v, 1, 0)$ , and thus it must also be adjacent to  $(2, w, 1, 0) = y_i^1$ . Since  $u$  is a neighbour of  $v$ ,  $(0, u, 3, 0)$  is a leaf in  $K_4T$  and hence is a leaf in  $S_2$ , thus  $(0, u, 3, 0) = x_l^2$ . Then  $(x_l^2, y_i^1) \in E(K_4T)$  and thus  $(z_l^1, z_i^2) \in E(F_T)$ . This means that given an edge  $(w, v) \in E(F(T))$  outside of  $F(T')$ , there exists an edge  $(z_l^1, z_i^2) \in E(F_T)$ . The other two cases are analogous.

Now consider an edge  $(z_i^1, z_j^2) \in F_T$  outside of  $F_{T'}$ . This means that  $(x_j^2, y_i^1) \in E(K_4T)$  with  $x_j^2 \in L_2$  and  $y_i^1 \in R_1$ . Since  $y_i^1 \in R_1$ , again by Corollary 2.6 we have that  $y_i^1 = (2, w, 1, 0)$  or  $y_i^1 = (2, w, 1, 0)$  if  $w$  is the root vertex of  $F(T)$ . Assume that  $y_i^1 = (2, w, 1, 0)$ , then  $x_j^2 = (0, v, 3, 0)$  or  $(0, v, 1, 2)$  and  $v$  is the unique neighbour of the leaf  $w$ , thus  $(v, w) \in F(T)$ . If  $y_i^1 = (2, w, 1, 0)$  with  $w$  the root of  $F(T)$ , we have that  $x_j^2 = (0, v, 3, 0)$  where  $v$  is the unique neighbour of  $w$ . Hence for every edge in  $F_T$  outside of  $F_{T'}$  there exists an edge in  $F(T)$  outside of  $F(T')$ . Thus  $F(T) = F_T$ .  $\square$

### 3.2 The general case

Let  $RK_4T$  be the graph obtained from  $K_4T$  as follows. For every set of vertices  $\{v_1, \dots, v_m\}$  such that  $N(v_i) = N(v_j)$  for  $i \neq j$ , we remove from  $K_4T$  the vertices  $\{v_2, \dots, v_m\}$ . Thus in  $RK_4T$  there are no two vertices having the same neighbourhood. Notice that the vertices of the graph  $RK_4T$  can also be separated into  $m_k$  sets  $S_1, \dots, S_k$  such that each  $S_i \cong KP_{m_i}$  for some  $m_i$ . Assume again that  $\{x_1^i, \dots, x_{m_i}^i\}$  and  $\{y_1^i, \dots, y_{m_i}^i\}$  denote the set of leaves and non-leaf vertices of  $S_i$  for  $1 \leq i \leq k$  respectively. We then have the following definition.

**Definition 3.6.** Assume  $V(RK_4T) = \cup_{i=1}^k V(S_i)$  with  $|V(S_i)| = m_i$ . The tree  $RF_T$  is defined as follows.

- The vertices of  $RF_T$  are  $V(F_T) = z_0 \cup (\cup_{i=1}^n \cup_{j=1}^{m_i} z_j^i)$ , and

- $N(z_0) = \cup_{j=1}^{m_n} z_j^n$ ,
- $(z_j^i, z_l^k) \in E(RF_T)$  if  $k = \min\{m > i : (x_l^m, y_j^i) \in E(RK_4T)\}$ .

The proof of the following Corollary is analogous to the proof of Lemma 3.5.

**Corollary 3.7.** *The trees  $RF_T$  and  $F(T)$  are isomorphic.*

Since  $T$  is not binary it is not fully determined by  $F(T)$ . Consider the set of vertices  $LR = \{x_1, \dots, x_{m_1}\}$  of  $K_4T$ , which have degree one in  $RK_4T$ . Assume that  $x_i$  has degree  $d_i$  in  $K_4T$  for  $1 \leq i \leq m_1$ . Then  $d_i = \frac{1}{2}(k_i - 1)(k_i - 2)$  for some integer  $k_i$ .

**Proposition 3.8.** *If a leaf  $x \in RK_4T$  has degree  $d = \frac{1}{2}(k - 1)(k - 2)$  in  $K_4T$  for some  $k$ . Then the corresponding leaf  $z \in F(T)$  has degree  $k$  in  $T$ .*

*Proof.* Let  $z$  be a leaf in  $F(T)$ , and let  $w$  be its unique neighbour in  $F(T)$ . Assume also that  $z$  lies on  $w$  direction  $i$ . If  $z \neq \star$ , then  $i \neq 0$  and consider a vertex  $v = (0, w, p, q) \in K_4T$  having 2 or 3 as the  $i$ th entry of  $p$  or as the  $i - l(p)$ th entry of  $q$ . They all have the same neighbourhood consisting of vertices of the form  $(2, z, p, 0)$ , and by Proposition 2.3, there are  $d = \frac{1}{2}(k - 1)(k - 2)$  such vertices, where  $d(z) = k$ . Thus all the vertices of the form of  $v$  correspond to a leaf  $x \in RK_4T$  which has degree  $d = \frac{1}{2}(k - 1)(k - 2)$ . If  $z = \star$ , we consider vertices of the form  $v = (0, z, p, q)$  having 2 or 3 as the  $i$ th entry of  $p$  or as the  $i - l(p)$ th entry of  $q$ , and they all have the same neighbourhood consisting of vertices of the form  $(2, w, p, 0)$ . Thus these last vertices correspond to a leaf  $x \in RK_4T$  which has the desired degree in  $K_4T$ . □

This means that we have fully recovered  $T$  from the graph  $K_4T$ .

## Declarations

On behalf of all authors, the corresponding author states that there is no conflict of interest. This manuscript has no associated data.

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